

# How to decrypt without keys with GlobalPlatform SCP02 protocol

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# SCP02

## Context

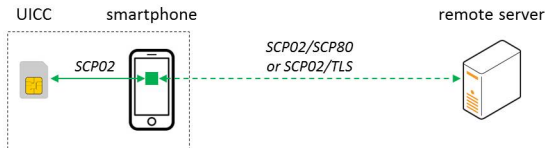
- Security protocol promoted by GlobalPlatform
- Element of a set of security protocols
  - “Symmetric” based: SCP02, SCP03, SCP80
  - “Asymmetric” based: SCP81, SCP10, SCP11
- Likely the most widely used SCP protocol

## Cryptographic functions

- Based on DES/3DES (encryption and MAC; cf. [ISO9797-1] and [ISO10116])

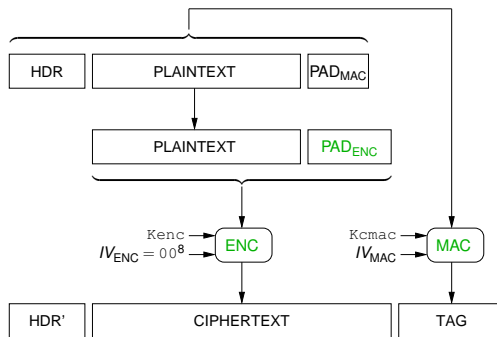
## Purpose

- Secure channel between an “*off card entity*” and a card
- Different security levels: integrity, confidentiality, both
- Remote card management (e.g., applet upload)



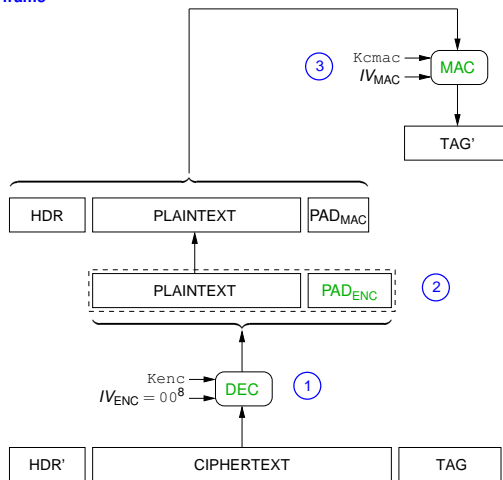
# Encryption and MAC

## MAC and encryption of a command frame



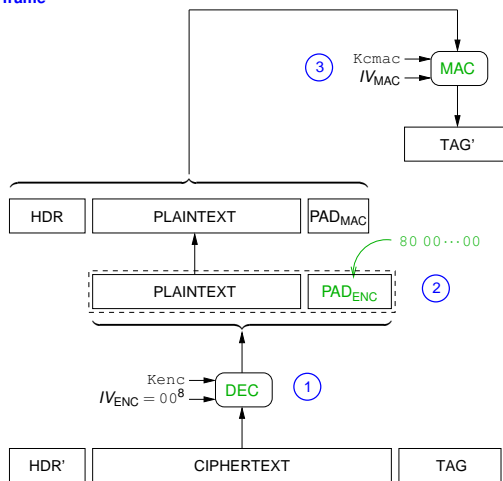
# Encryption and MAC

## Decryption of a command frame



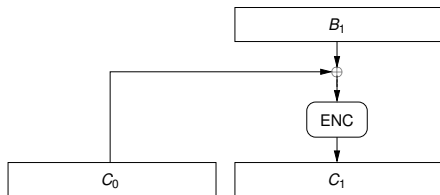
# Encryption and MAC

## Decryption of a command frame



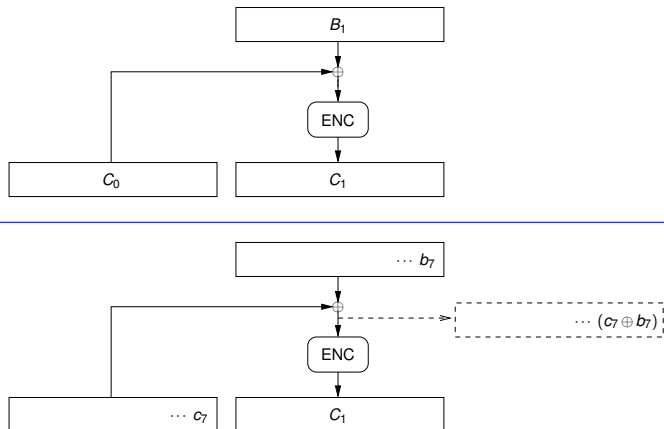
# Encryption and MAC

## CBC encryption



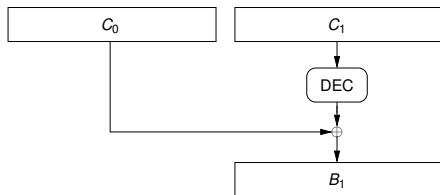
# Encryption and MAC

## CBC encryption



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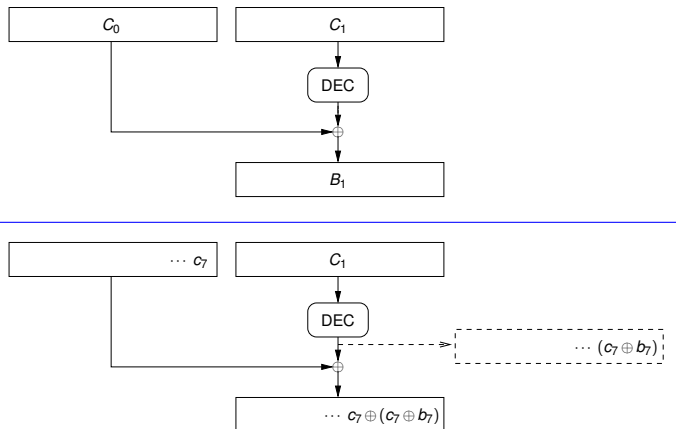
## CBC decryption





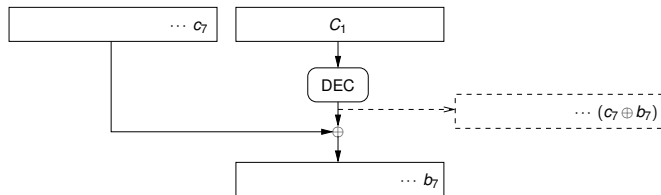
# Encryption and MAC

## CBC decryption



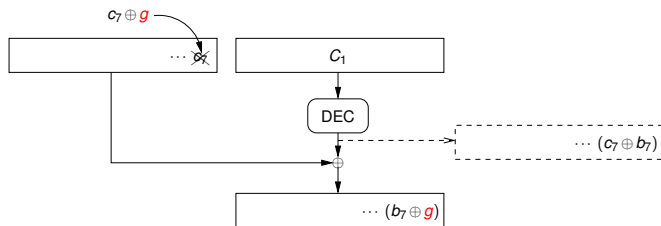
# Encryption and MAC

## CBC decryption and malleability



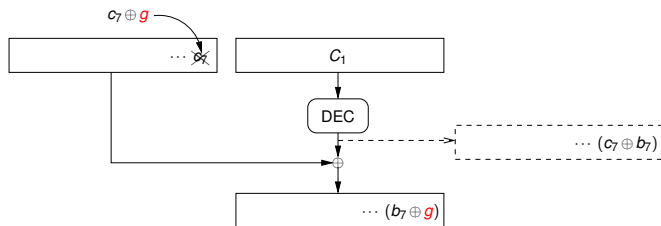
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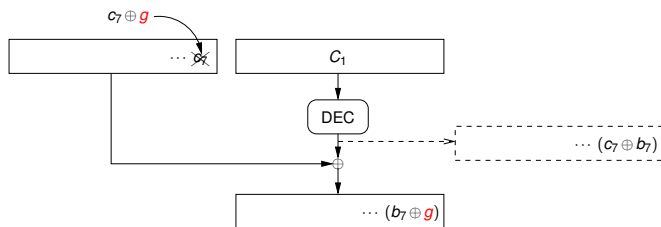


$g = 13 \Rightarrow$

$\Rightarrow$  invalid padding

# Encryption and MAC

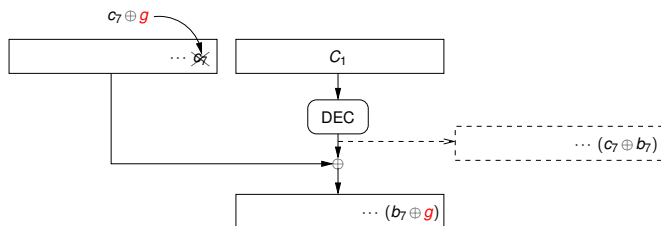
## CBC decryption and malleability



$g = 13 \Rightarrow$   $\Rightarrow$  invalid padding  
 $g = 14 \Rightarrow$   $\Rightarrow$  invalid padding

# Encryption and MAC

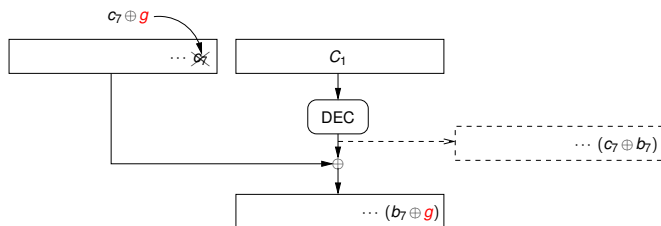
## CBC decryption and malleability



$g = 13 \Rightarrow \Rightarrow$  invalid padding  
 $g = 14 \Rightarrow \Rightarrow$  invalid padding  
 $g = 15 \Rightarrow \Rightarrow$  invalid padding

# Encryption and MAC

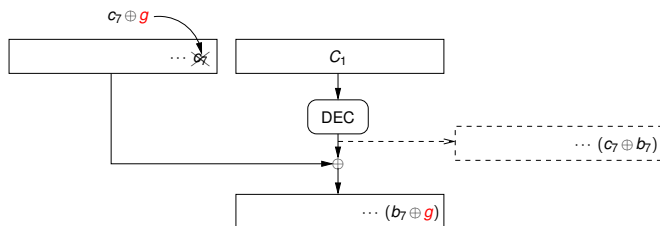
## CBC decryption and malleability



$g = 13$	$\Rightarrow$	$\Rightarrow$	invalid padding
$g = 14$	$\Rightarrow$	$\Rightarrow$	invalid padding
$g = 15$	$\Rightarrow$	$\Rightarrow$	invalid padding
$g = 16$	$\Rightarrow$	$\Rightarrow$	invalid padding

# Encryption and MAC

## CBC decryption and malleability

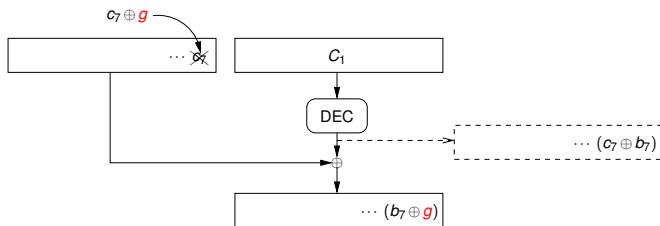


$g = 13$	$\Rightarrow$	$\Rightarrow$	invalid padding
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$g = 16$	$\Rightarrow$	$\Rightarrow$	invalid padding
$g = 17$	$\Rightarrow$	$\Rightarrow$	valid padding



# Encryption and MAC

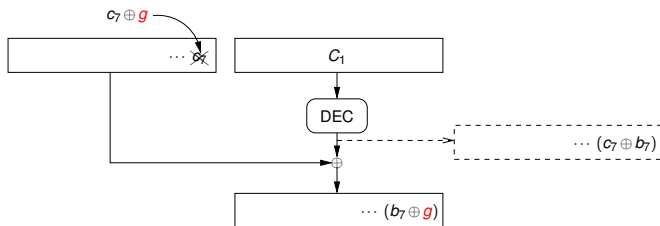
## CBC decryption and malleability



$g = 13$	$\Rightarrow$	$b_7 \oplus g = 76$	$\Rightarrow$	invalid padding	
$g = 14$	$\Rightarrow$	$b_7 \oplus g = 77$	$\Rightarrow$	invalid padding	
$g = 15$	$\Rightarrow$	$b_7 \oplus g = 78$	$\Rightarrow$	invalid padding	
$g = 16$	$\Rightarrow$	$b_7 \oplus g = 79$	$\Rightarrow$	invalid padding	
$g = 17$	$\Rightarrow$	$b_7 \oplus g = 80$	$\Rightarrow$	valid padding	$\Rightarrow b_7 = g \oplus 80 = 97$

# Encryption and MAC

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- The **validity of padding data** indicates whether  $b_7$  can be found or not.
- Technique called “padding oracle attack” due to Vaudenay in 2002 [V02].

# Generic padding oracle attack

## Padding oracle

- The remote server wants to send secret data  $B_1 = S$  to the card.
- The remote server embeds  $B_1 = S$  in an encrypted SCP02 command  $C_0 \| C_1 \| C_2$ .
- The attacker  $\mathcal{A}$  gets  $C_0 \| C_1 \| C_2$ , replaces it with  $\tilde{C}_0 \| C_1$  (adding  $g$ ), and sends the result to the card.

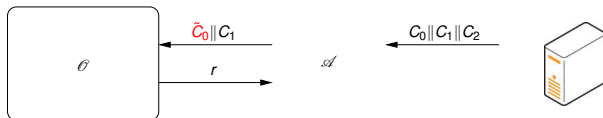


- if  $r$  = "valid padding" then  $\mathcal{A}$  **decrypts one secret byte**
- if  $r$  = "invalid padding" then  $\mathcal{A}$ 
  - picks another value  $g$
  - goes to a.

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# Generic padding oracle attack

## Building the padding oracle $\mathcal{O}$

- Operations during decryption process

(A)	(B)
1. decryption	1. decryption
2. padding data: <b>invalid</b>	2. padding data: <b>valid</b>
3. <del>MAC</del>	3. MAC

# Generic padding oracle attack

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- Oracle  $\mathcal{O}$  based on error message (e.g., WTLS, see [V02])

ERR_DEC	⇒	invalid padding data	(A)
ERR_MAC	⇒	valid padding data	(B)

# Generic padding oracle attack

## Building the padding oracle $\mathcal{O}$

- Operations during decryption process

(A)	(B)
1. decryption	1. decryption
2. padding data: <b>invalid</b>	2. padding data: <b>valid</b>
3. <del>MAC</del>	3. MAC

- Oracle  $\mathcal{O}$  based on error message (e.g., WTLS, see [V02])

ERR_DEC	⇒	invalid padding data	(A)
ERR_MAC	⇒	valid padding data	(B)

- Oracle  $\mathcal{O}$  based on computation time (e.g., TLS 1.0, see [CHVV03])

time ↘	⇒	<del>MAC</del>	⇒	invalid padding data	(A)
time ↗	⇒	MAC	⇒	valid padding data	(B)

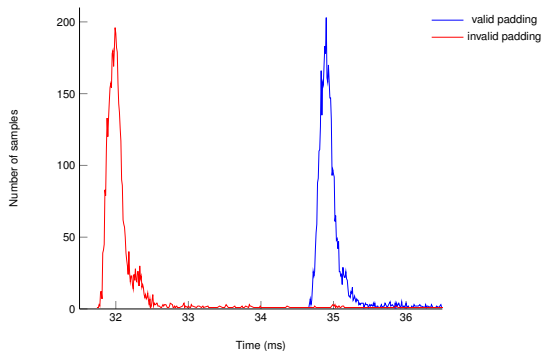
# Padding oracle in SCP02-compliant smart cards

- The smart card sends *always* a response (status word).
- Invalid padding data or invalid MAC  $\Rightarrow$  same error code



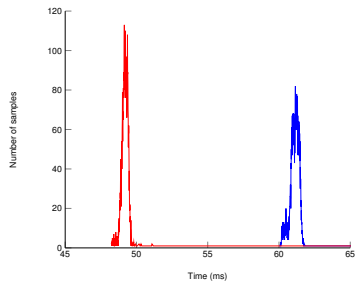
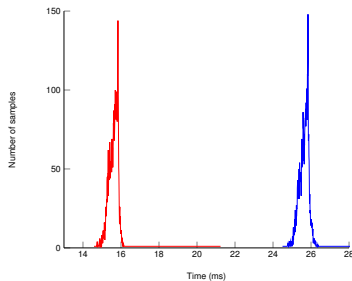
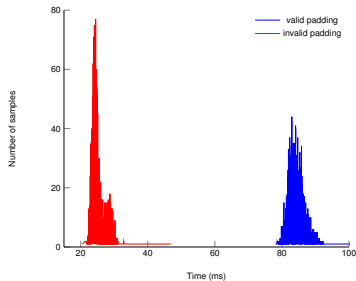
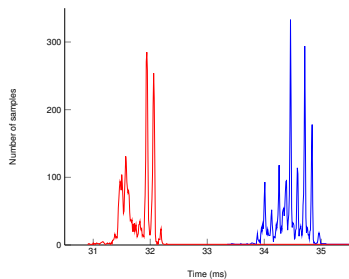
# Padding oracle in SCP02-compliant smart cards

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- Invalid padding data or invalid MAC  $\Rightarrow$  same error code



- The **card response time** reflects the card computation time  $\Rightarrow$  suitable padding oracle  $\mathcal{O}$

# Padding oracle in SCP02-compliant smart cards



# Padding oracle attack against SCP02-compliant smart cards

## Practical experiments

- Experimental setting: card connected to a card reader (4 card readers, wired and wireless)
- 10 smart cards from 6 card manufacturers
- SIM cards, generic Java cards
  
- Experiment: find a 16-byte secret key sent to the smart card in an encrypted SCP02 command
- 300 experiments/card  $\Rightarrow$  100 % success
- Practical complexity  $\in [127.75, 133.38]$  close to best average case (128)
- Time to find 16 bytes: 2.7 mn to 11.4 mn (variable response time from the smart card)

M	C	$\mu_W$ (ms)	$\mu_R$ (ms)	$t_{min}$ (ms)	$m$	$\tau_+$ (%)	$K_W$	$K_R$	Z	Z/n
1	A	39.60	42.59	41.00	28	0.16	1	3	2055.71	128.48
	B	40.19	43.94	42.00	28	0.44	1	3	2077.78	129.86
2	C	25.17	84.34	75.00	0	0.00	1	2	2043.95	127.75
	D	26.64	34.36	32.00	0	0.00	1	2	2066.54	129.16
3	E	15.61	25.65	23.00	0	0.00	1	2	2134.03	133.38
4	F	31.81	34.48	33.00	28	0.48	1	3	2109.71	131.86
	G	15.64	18.53	17.00	0	0.28	1	3	2103.62	131.48
5	H	25.18	84.86	72.00	0	0.00	1	2	2048.34	128.02
6	I	25.90	35.85	32.00	0	0.06	1	3	2108.60	131.79
	J	14.32	19.92	17.50	0	0.10	1	2	2094.85	130.93

# Padding oracle attack against SCP02-compliant smart cards

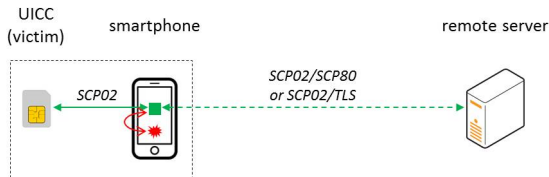
## Requirements



1. The attacker sits **between the remote server and the card** at a point where she can directly eavesdrop on SCP02 encrypted commands and send modified commands to the card.
2. The attacker is able to **discriminate response times** corresponding to a valid and an invalid padding.
3. The remote server **repeatedly** sets up a (new) secure channel with the card.
4. The **same secret information** is sent through each such secure channel.
5. The secret information is sent at a **predictable position**.

NB: req. 4  $\Rightarrow$  req. 3 (and 5)

# Attack scenario

## Trojan in a smartphone



1. The victim downloads from a popular store an infected application (embedding a Trojan ) into his smartphone.
2. The Trojan gets access to the memory space of the legitimate application  (through privileges escalation).
3. The Trojan can apply the attack: it reads, and modifies the encrypted SCP02 commands received by the legitimate application.

# Countermeasures for SCP02-compliant smart cards

- Correct implementation (not possible for deployed cards)
- Use additional security mechanisms
- Use `PUT_KEY` command to send sensitive data
- Do not send too many times the same data (server side)

# Responsible disclosure

## Disclosure (October 2017-April 2018)

- Card manufacturers
- GlobalPlatform
  
- ANSSI has been also informed.

## Publication

- Paper accepted at conference [CHES 2018](#)  
*Attacking GlobalPlatform SCP02-compliant Smart Cards Using a Padding Oracle Attack*,  
Gildas Avoine (INSA Rennes), Loïc Ferreira
- Published in IACR Transactions on Cryptographic Hardware and Embedded Systems [AF18].  
To be presented at CHES conference in [September](#) 2018.

# Take-away

- The padding oracle attack against SCP02-compliant smart cards is possible because of
  - a [theoretical flaw](#) lying in the SCP02 protocol (Encrypt-and-MAC scheme),
  - exploited by means of a [timing side-channel](#) provided by the smart cards (implementation).
- Several [requirements](#) to be fulfilled in order for the attack to be successful.
- Practical attack
  - Experimental setting: 10 smart cards from 6 manufacturers.
  - [How many](#) smart cards impacted in real life?
- SCP02 is now [deprecated](#) (March 2018): use [SCP03](#) instead.



Thank you

Questions?



# References

- [SCP02] GlobalPlatform. *GlobalPlatform – Card Specification*, version 2.3.1, ref. GPC\_SPE\_034, March 2018.
- [SCP03] GlobalPlatform. *GlobalPlatform Card Technology – Secure Channel Protocol '03' – Card Specification v2.2 – Amendment D*, version 1.1, ref. GPC\_SPE\_014, July 2014.
- [ISO9797-1] ISO/IEC JTC 1/SC 27. *ISO/IEC 9797-1:2011 – Information technology – Security techniques – Message Authentication Codes (MACs) – Part 1: Mechanisms using a block cipher*, 2011.
- [ISO10116] ISO/IEC JTC 1/SC 27. *ISO/IEC 10116:2017 – Information technology – Security techniques – Modes of operation for an n-bit block cipher*, 2017.
- [ISO7816-4] ISO/IEC JTC 1/SC 17. *ISO/IEC 7816-4:2013 – Information technology – Identification cards – Integrated circuit cards – Part 4: Organization, security and commands for interchange*, 2013.
- [CHVV03] B. Canvel, A. Hiltgen, S. Vaudenay, M. Vuagnoux. *Password interception in a SSL/TLS channel*. In D. Boneh, editor, *Advances in Cryptology – CRYPTO 2003*. LNCS, vol. 2729, pp. 583-599. Springer, 2003.
- [V02] S. Vaudenay. *Security Flaws Induced by CBC Padding – Applications to SSL, IPSEC, WTLS...* In L. R. Knudsen, editor, *Advances in Cryptology – EUROCRYPT 2002*. LNCS, vol. 2332, pp. 534-545. Springer, 2002.
- [ST16] M. Sabt, J. Traoré. *Cryptanalysis of GlobalPlatform Secure Channel Protocols*. In L. Chen, D. McGrew, C. Mitchell, editors, *Security Standardisation Research – SSR 2016*. LNCS, vol. 10074, pp. 62-91. Springer, 2016. <https://eprint.iacr.org/2017/032>
- [AF18] G. Avoine, L. Ferreira. *Attacking GlobalPlatform SCP02-compliant Smart Cards Using a Padding Oracle Attack*. In IACR Transactions on Cryptographic Hardware and Embedded Systems, Issue 2, pp. 149-170, May 2018. <https://tches.iacr.org/index.php/TCHES/article/view/878>